

# Exercises on Randomized Complexity

## CSCI 6114 Fall 2021

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Recall that we have a language  $L$  is in...

	BPP	RP	coRP
$x \in L \Rightarrow$	$Pr[A(x) \text{ accepts}] \geq 2/3$	$Pr[A(x) \text{ accepts}] \geq 2/3$	$Pr[A(x) \text{ accepts}] = 1$
$x \notin L \Rightarrow$	$Pr[A(x) \text{ accepts}] \leq 1/3$	$Pr[A(x) \text{ accepts}] = 0$	$Pr[A(x) \text{ accepts}] \leq 1/3$

Last time we saw that POLYNOMIAL IDENTITY TESTING is in coRP.

1. Show that  $P^{BPP} = BPP$ , and then show that  $BPP^{BPP} = BPP$ .
2. Show that  $coBPP = BPP$ .
3. Give an alternative characterization of BPP similar to the verifier definition of NP. Here, instead of a “witness”, think of the Verifier  $V(x, r)$  as taking in the input  $x$  and a random string  $r$ .
4. Show that  $RP \subseteq NP$ . Show that  $NP \subseteq BPP$  iff  $NP = RP$ . *Hint:* BPP and RP are both closed under  $\leq_m^p$ , so NP is contained in one of these classes iff SAT is.

We can now update our table above to:

$Pr(A(x) \text{ accepts})$	BPP	RP	coRP	NP	coNP	PP
$x \in L \Rightarrow$	$\geq 2/3$	$\geq 2/3$	$= 1$	$> 0$	$= 0$	$> 1/2$
$x \notin L \Rightarrow$	$\leq 1/3$	$= 0$	$\leq 1/3$	$= 0$	$> 0$	$\leq 1/2$

5. (a) Let  $L \in BPP$ . Let  $L'$  be the poly( $n$ )-concatenation of  $L$ , that is, a tuple  $(x_1, \dots, x_n) \in L'$  iff all  $x_i \in L$ . Show that  $L' \in BPP$ .
- (b) Show that  $NP^{BPP} = NP^{BPP[1]}$  where the latter means the oracle is queried only once. *Hint:* Use nondeterminism to guess the oracle answers, and use the one query at the end to verify the guesses.
- (c) Show that  $NP^{BPP} \subseteq BPP^{NP}$ .
- (d) Use the preceding to show that if  $NP \subseteq BPP$  then  $PH \subseteq BPP$ . (We’ll see next week that  $BPP \subseteq \Sigma_2P \cap \Pi_2P$ , so in fact the latter implies that PH collapses.)

## Resources

- TODO